

## Varieties

a syntax semantics duality.

- (1) We saw that a category with products  $\mathcal{C}$  can be seen as an "operational" theory, and this is even more evident when we specify a b.o. functor  $\text{Fin}/X \xrightarrow{\text{op } \ell} \mathcal{C}$  preserving products. Indeed,

$$\text{Mod}(\mathcal{C}) \xrightarrow{\ell^*} \text{Prod}(\text{Fin}/X, \text{Set})$$

$$\downarrow \text{is } X$$

$$\text{Set}$$

the functor  $\ell^*$  gives us a forgetful functor that delivers a family of relevant sorts for the "theory" axiomatized by  $\mathcal{C}$ .

- (2) We saw that the inclusion

$$\text{Mod}(\mathcal{C}) \xrightarrow{i} \text{Set}^{\mathcal{C}}$$

$\begin{array}{ccc} & \xleftarrow{L} & \\ & \perp & \\ & \xrightarrow{i} & \end{array}$

- creates limits (trivially).
- creates directed & sifted colimits (they commute with products in  $\text{Set}$ ).
- has a left adjoint (only sketched, postponed to orthogonality theory).
- has a dense subcategory  $\mathcal{C}^{\text{op}} \xrightarrow{\ell} \text{Mod}(\mathcal{C})$ .

- Q1: But can we characterize "algebraic categories" ?
- Q2: Do the properties we listed characterize them ?

So, the honest answer to Q1 is yes but in many equivalent ways. In this lecture we do not choose the original one provided by Lawvere, nor the most essential one, which emerged in a long process of maturation involving several people (Pedicchio, Linton, Vitale, Adamek, Rosicky, ...). We choose the one that most easily generalises to other contexts.

Let us do the first steps in the direction of answering the question.

Remark. In general the objects of the form  $f c$

$$f: C \rightarrow \text{Psh}(C)$$

are tiny, i.e. the hom-functor preserves all colimits, indeed

~~$$\text{colim}_i (f c_i) \cong f(\text{colim}_i c_i)$$~~

~~$$\text{Psh}(C)(f c, \text{colim}_i P_i) \cong (\text{colim}_i P_i)(c)$$~~

$$\text{colimts are pointwise} \cong \text{colim} [P_i(c)]$$

$$\text{Yoneda} \cong \text{colim Psh}(C)(f c, P_i)$$

Now, because the inclusion creates sifted colimits,

(3)

$$\text{Mod}(C) \xrightarrow{i} \text{Set}^C$$

it follows that the representable models are "tiny" with respect to those.

$$\text{Mod}(C)(f_C, \text{colim } x_i) \stackrel{\text{Fullman}}{\cong} \text{Set}^C(f_C, i(\text{colim } x_i))$$

$$i \text{ creates sifted colimits} \cong \text{Set}^C(f_C, \text{colim } x_i)$$

$$\boxed{\text{sifted} = \text{directed} + \text{reflexive coeq}}$$

□.

Def Finitely presentable object

Examples: ... finite stuff.

Reflexive equalizers & projectivity

Examples: free algebras.

So, in a "variety" whatever this should be we have dense generator made of these

finite, projective objects.

Prop the full subcategory of these objects is closed under coproducts

(multisorted)

Def A variety is a complete category with a dense generator made of "finite projective" objects.

Morphisms of varieties -

- $\text{Fin}^{\text{op}}$ , after all, is a "theory" itself, and its models are sets -

$$\text{Mod}(\text{Fin}^{\text{op}}) \simeq \text{Set} -$$

The functor specifying the sets of a theory  $\ell: \text{Fin}^{\text{op}} \rightarrow \mathcal{C}$  induces a functor

$$\text{Mod}(\mathcal{C}) \xrightarrow{\ell^*} \text{Mod}(\text{Fin}^{\text{op}})$$

- preserving limits
- preserving sifted colimits - (AFT  $\Rightarrow$  has a left adjoint!!)

Prop The left adjoint preserve finite projections.

Pf A morphism of varieties is a right adjoint preserving sifted colimits -

Ex  $\text{Set}^{\mathcal{C}}$  is a variety and also the inclusion

$$\text{Mod}(\mathcal{C}) \longrightarrow \text{Set}^{\mathcal{C}} \text{ is}$$

a morphism of varieties!

Rem Interpretations ...

# Duality

(5)



- Mod was essentially already defined.
- In the other direction we map  $V$  to the full subcategory of "finite projectives"  $V_{\omega}^P$

Rec (Makkai-Lovvorn)  $Mod \simeq Prod(-, Set)$

$V_{\omega}^P \simeq Ver(-, Set)$

Set is the dualizing object.